

**THE BEHAVIOURAL APPROACH
to
SYSTEMS and CONTROL**

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OUTLINE

- **Introduction**
- **The behavior**
- **Modeling interconnected systems**
- **Elimination**
- **Controllability**
- **Representations**
- **Algorithmic issues**
- **Control as interconnection**

SALIENT FEATURES

- **Dynamical system =
a behavior**

input/output structure =
important special case

- **First principles models
⇒ latent variables**

state variables =
important special case

- **Modeling complex systems
= tearing & zooming**

i/o cascade and feedback =
limited special case

- **Control = interconnection**

feedback =
important special case

SYSTEM

Let w_1, w_2, \dots, w_q be variables whose dynamic relation we wish to describe (modeling), or analyze, or design (control and synthesis)

.....

Ingredients:

‘Time set’ \mathbb{T}

(today, $\mathbb{T} = \mathbb{R}$ but framework covers \mathbb{Z} , DES, hybrid systems)

w_k takes on its value in \mathbb{W}_k .

Yields $\mathbb{W} = \mathbb{W}_1 \times \mathbb{W}_2 \times \dots \times \mathbb{W}_k$

‘signal space’

How do we express the ‘laws’?

SYSTEM (continued)

Classical framework:

some of w_k 's act as
inputs, causes, stimuli,
the other w_k 's act as
outputs, effects, responses.

system = I/O map
transfer function, etc.

internal initial conditions
→ familiar I/S/O representation

$$\frac{d}{dt}x = f(x, u), \quad y = h(x, u)$$

via output to input cascade + feed-
back interconnection : complex sys-
tems

very successful framework in signal
processing, control, simulation, etc.

SYSTEM (continued)

Limitations:

- cause/effect, stimulus/response is often simply not physical!
- very awkward framework for use in first principles modeling: what is the signal flow graph?
- not well-suited for computer assisted 'object oriented' modeling!!!
- unnecessarily limits model representations.

THE BEHAVIOR

The *behavioral approach* takes the feasible trajectories

$$w : \mathbb{T} \rightarrow \mathbb{W}$$

$$w = (w_1, w_2, \dots, w_q)$$

as the object of study.

Totality of feasible trajectories = \mathfrak{B}
the '*behavior*'; $\mathfrak{B} \subset \mathbb{W}^{\mathbb{T}}$

$w \in \mathfrak{B}$: the system allows
the trajectory $w = (w_1, w_2, \dots, w_q)$

$w \notin \mathfrak{B}$: the system forbids
the trajectory $w = (w_1, w_2, \dots, w_q)$

Note: in I/O-systems
 $\mathfrak{B} =$ all (u, y) -pairs.

Note: analogy with formal language:
some words OK, some words not OK.

THE BEHAVIOR (continued)

'Word' examples

- Planetary orbits
- Port behavior of an electrical circuit

- Force/position behavior in mechanical systems

- other frameworks, that have yielded away from I/O modeling:
 - across/through variables
(bondgraphs)
 - intensive/extensive variables
(thermodynamics)
 - formal languages, PDE's
(cs, physics)

BEHAVIORAL EQUATIONS

Differential equations

$(T = \mathbb{R}, W = \mathbb{R}^w)$,

$$F(w(t), \frac{dw}{dt}(t), \dots, \frac{d^n w}{dt^n}(t)) = 0$$

model $\cong F$.

Linear case

$$R_0 w + R_1 \frac{dw}{dt}(t) + \dots + R_n \frac{d^n w}{dt^n}(t) = 0$$

R_0, R_1, \dots, R_n : constant matrices =
model parameters.

behavior = all solutions

BUT! Models are seldomly given this
way!

Auxiliary variables !!!

LATENT VARIABLES

Behavioral equations:

$$F\left(w(t), \frac{dw}{dt}(t), \dots, \frac{d^n w}{dt^n}(t), \right. \\ \left. \ell(t), \frac{d\ell}{dt}(t), \dots, \frac{d^n \ell}{dt^n}(t)\right) = 0$$

Behavior: all ‘solutions’ $w : \mathbb{R} \rightarrow \mathbb{W}$
that is, all $w : \mathbb{R} \rightarrow \mathbb{W}$, for which
there exists $\ell : \mathbb{R} \rightarrow \mathbb{L}$, such that
 (w, ℓ) is a solution.

w : ‘manifest’ variables

ℓ : ‘latent’ variables

Example:

$$\frac{d}{dt}x = f(x, u), y = h(x, u) \\ w = (u, y)$$

x : latent, (u, y) : manifest.

Latent \cong internal,
but internal \neq state.

INTERCONNECTED SYSTEM

Graph with leaves

nodes = modules

edges = connections

= pairing of terminals

leaves = external terminals

Think about an electrical circuit

MODELING

A computer-assistance oriented procedure for obtaining a model for an interconnected system

Central notions involved:

- **terminals**
- **modules**
- **interconnection architecture**

.....

modules = building blocks

**terminals = links between subsystems
carry variables that
vary with time**

**interconnection = the way the subsystems
architecture are linked**

TERMINALS

A *terminal* is specified by its *type*.

The *type* implies an ordered set of *variables*.

Examples

Type of terminal	Variables
electrical	(voltage, current)
mechanical (1-D)	(force, position)
mechanical (2-D)	(force, torque, position, attitude)
thermal	(temp, heat flow)
fluidic	(pressure, flow)
m-dim input	(u_1, u_2, \dots, u_m)
p-dim output	(y_1, y_2, \dots, y_p)
etc.	etc.

MODULES

A *module* is specified by *its type*,
its representation and
its parameter values.

The *type* specifies an ordered set of
terminals (t_1, t_2, \dots, t_n)

This specifies an ordered set of vari-
ables (w_1, w_2, \dots, w_n)

The *representation* specifies
the behavior of these variables.

This representation contains param-
eters. The *parameter values* specify
their values.

By specifying a module we obtain
the behavior of the variables
 (w_1, w_2, \dots, w_n)
on the terminals of the module.

MODULES (continued)

Examples

Module	Repr.	Par. value
resistor	imp	R in ohms
resistor	adm	G in mhos
cap	default	C in farad
Δ	default	R in ohms
transformer	default	turns ratio
m-port Imp	tf. fn.	$G \in \mathbb{R}^{m \times m}(\xi)$
m-port imp	state repr.	(A, B, C, D)
m-port imp	kernel repr.	$R \in \mathbb{R}^{m \times 2m}[\xi]$
m-port imp	im repr.	$M \in \mathbb{R}^{2m \times m}[\xi]$
mass	default	m in kgr
pendulum	default	m and L
2 inlet tank	default	geometry
(m,p) lin sys	state repr.	(A,B,C,D)
etc.	etc.	etc.

MODULES (continued)

terminals = (t_1, t_2)

t_1, t_2 : both electrical

$t_1 \rightarrow (V_1, I_1), t_2 \rightarrow (V_2, I_2)$

behavioral equations:

$$V_1 - V - 2 = RI_1, I_1 = I_2$$

.....

terminals= (t_1, t_2)

behavioral equations

$$C \frac{d}{dt}(V_1 - V_2) = I_1, I_1 = I_2$$

.....

terminals= (t_1, t_2, t_3)

latent variables: I'_1, I'_2, I'_3

behavioral equations:

$$V_1 - V_2 = RI'_3, V_2 - V_3 = RI'_1, V_3 - V_1 = RI'_2$$

$$I_1 = I'_3 - I'_2, I_2 = I'_1 - I'_3, I_3 = I'_2 - I'_1$$

MODULES (continued)

terminals= $(t_1, t_2, \dots, t_m, \text{ground})$

all electrical

latent variables: x

behavioral equations:

$$\frac{d}{dt}x = Ax + BI', \quad V' = Cx + DI'$$

$$(I_1, \dots, I_m, I_{m+1}) = (I'_1, \dots, I'_m, -I'_1 - I'_2 \dots - I'_m)$$

$$(V_1, \dots, V_m, V_{m+1}) = (V'_1 + V, \dots, V'_m + V, V)$$

INTERCONNECTION ARCHITECTURE

The interconnection architecture is
a list of
terminal pairs

Connect	
t'	t''
t'''	t''''
etc.	etc.
etc.	etc.

$\{t', t''\}$ can be such a pair only if the
type of t' is suitably adapted to the
type of t''

“adapted”

→ same type (electrical, mechanical, thermal)

→ output to input for ‘logical’
connections.

INTERCONNECTION ARCHITECTURE (continued)

The interconnection architecture imposes restrictions on the variables 'living' on the associated terminals.

Electrical: $t' \rightarrow V', I', t' \rightarrow V'', I''$
restriction: $V' = V'', I' + I'' = 0$

1-D mechanical:

$t' \rightarrow F', q', t'' \rightarrow F'', q''$
restriction: $F' + F'' = 0, q' = q''$

2-D mechanical:

$t' \rightarrow x', y', \theta', X', Y', T'$
 $t'' \rightarrow x'', y'', \theta'', X'', Y'', T''$
restriction: $x' = x'', y' = y'', \theta' = -\theta'', X' + X'' = 0, Y' + Y'' = 0, T' = T''$

thermal $t' \rightarrow T', Q', t' \rightarrow T'', Q''$
restriction: $T' = T'', Q' + Q'' = 0$

logical $t' \rightarrow u', t' \rightarrow y''$
restriction: $y = u$

MODEL GENERATION

So in order to obtain a model specify:

- Modules M_1, M_2, \dots, M_N
type + representation +
parameter values.
This yields a list of terminals t_1, t_2, \dots, t_N
and a behavior \mathfrak{B}' for the variables living on the terminals.
- Interconnection architecture on
 t_1, t_2, \dots, t_N
this yields a behavior \mathfrak{B}'' for the variables living on the terminals
- $\mathfrak{B}' \cap \mathfrak{B}''$ = the behavior of the interconnected system
contains latent variables and manifest variables.
- Elimination of latent variables \rightarrow
 \mathfrak{B}

I/O and INTERCONNECTIONS

Consider 2 tanks:

$$\begin{array}{l|l} \frac{d}{dt}h_1 = F_1(h_1, p_1, p_2) & \frac{d}{dt}h_2 = F_3(h_3, p_3, p_4) \\ f_1 = H_1(h_1, p_1) & f_3 = H_3(h_3, p_3) \\ f_2 = H_2(h_1, p_2) & f_4 = H_4(h_3, p_4) \\ \text{input: } p_1, p_2 & \text{input: } p_3, p_4 \\ \text{output: } f_1, f_2 & \text{output: } f_3, f_4 \end{array}$$

$$\text{Interconnection: } p_2 = p_3, f_2 + f_3 = 0$$

input=input; output=output

very many such examples (e.g., in mechanics, heat transfer)

RLC circuit

! Model the relation between V_1, I_1, V_2, I_2 .

RLC circuit (continued)

Module	type	terminals	par. value
I	connector	1,3,4	
II	capacitor	5,6	C
III	resistor	7,8	R_C
IV	resistor	9,10	R_L
V	inductor	11,12	L
VI	connector	13,14,2	

Architecture: connect terminals

3	5
4	7
6	9
8	11
10	13
12	14

RLC circuit (continued)

Each terminal has two real variables:
voltage V , current I .

Generate the behavioral equations
via module laws & interconnections:

Modules	
$I_1 + I_3 + I_4 = 0$	$V_1 = V_3 = V_4$
$I_2 + I_{13} + I_{14} = 0$	$V_2 = V_{13} = V_{14}$
$I_5 + I_6 = 0$	$C \frac{d}{dt}(V_5 - V_6) = I_5$
$I_7 + I_8 = 0$	$V_7 - V_8 = R_C I_7$
$I_9 + I_{10} = 0$	$V_9 - V_{10} = R_L I_9$
$I_{11} + I_{12} = 0$	$V_{11} - V_{12} = L \frac{d}{dt} I_{11}$

Interconnections	
$V_3 = V_5$	$I_3 = -I_5$
$V_4 = V_7$	$I_4 = -I_7$
$V_6 = V_9$	$I_6 = -I_9$
$V_8 = V_{11}$	$I_8 = -I_{11}$
$V_{10} = V_{13}$	$I_{10} = -I_{13}$
$V_{12} = V_{14}$	$I_{12} = -I_{14}$

Model for V_1, I_1, V_2, I_2 : manifest,
 $V_3, I_3, \dots, V_{14}, I_{14}$: latent.

RLC circuit (continued)

! Eliminate $V_3, I_3, \dots, V_{14}, I_{14}$!

After suitable manipulations, we obtain:

CART with DOUBLE PENDULUM

! Model relation between u and y .

DOUBLE PENDULUM CART (continued)

View as interconnection of
5 modules

M_1 : pointmass with 2 rigid pins and
1 hinge
terminals (t_1, t_2, t_3)
2-D mechanical

M_2 : pendulum with rigid bar
and 1 hinge
terminals (t_4, t_5)
2-D mechanical

M_3 : pendulum with rigid base
terminal (t_6)
2-D mechanical

M_4 : horizontal drive
terminal (t_7)
2-D mechanical

M_5 : horizontal free motion pin
terminal (t_8)
2-D mechanical

DOUBLE PENDULUM CART (continued)

Interconnection architecture

Connect	
t_1	t_7
t_2	t_8
t_3	t_4
t_5	t_6

DOUBLE PENDULUM CART (continued)

Description of modules

M_1 : pointmass with 2 rigid pins
and 1 hinge

parameter values mass, $\Delta\theta$

variables:

$x_1, y_1, \theta_1, X_1, Y_1, T_1$

$x_2, y_2, \theta_2, X_2, Y_2, T_2$

$x_3, y_3, \theta_3, X_3, Y_3, T_3$

latent variables:

x, y, θ, X, Y

DOUBLE PENDULUM CART (continued)

Behavioral equations:

$$x_1 = x_2 = x_3 = x$$

$$y_1 = y_2 = y_3 = y$$

$$\theta_2 = \theta_1 + \underline{\Delta\theta}$$

$$X = X_1 + X_2 + X_3$$

$$Y = Y_1 + Y_2 + Y_4$$

$$\underline{\text{mass}} \frac{d^2x}{dt^2} = X$$

$$\underline{\text{mass}} \frac{d^2y}{dt^2} = Y$$

$$T_1 + T_2 = 0$$

$$T_3 = 0$$

DOUBLE PENDULUM CART (continued)

M_2 : pendulum with rigid base and one hinge

parameter values length, mass

variables:

$$x_1, y_1, \theta_1, X_1, Y_1, T_1$$

$$x_2, y_2, \theta_2, X_2, Y_2, T_2$$

latent variables:

$$x, y, \theta, X, Y$$

DOUBLE PENDULUM CART (continued)

Behavioral equations:

$$x_1 = x - \text{length} \cos \theta_1$$

$$x_2 = x$$

$$y_1 = y - \text{length} \sin \theta_1$$

$$y_2 = y$$

$$X = X_1 + X_2$$

$$Y = Y_1 + Y_2$$

$$\text{mass} \frac{d^2 x}{dt^2} = X$$

$$\text{mass} \frac{d^2 y}{dt^2} = Y - \text{mass} g$$

$$\text{mass} (\text{length})^2 \frac{d^2 \theta}{dt^2} = T_1 - \text{mass} g \text{length} \cos \theta_1$$

$$T_2 = 0$$

DOUBLE PENDULUM CART (continued)

M_3 : Pendulum with rigid base
parameter values: length, mass

variables:

$$x, y, \theta, X, Y, T,$$

behavioral equations:

$$\begin{aligned}\underline{\text{mass}} \frac{d^2 x}{dt^2} &= X \\ \underline{\text{mass}} \frac{d^2 y}{dt^2} &= Y - \underline{\text{mass}} g \\ \underline{\text{mass}} (\underline{\text{length}})^2 \frac{d^2 \theta}{dt^2} &= T - \underline{\text{mass}} g \underline{\text{length}} \cos \theta\end{aligned}$$

DOUBLE PENDULUM CART (continued)

M_4 : horizontal drive

external force input $u : \mathbb{R} \rightarrow \mathbb{R}$

variables:

$$x, y, \theta, X, Y, T$$

behavioral equations:

$$y = 0$$

$$\theta = 0$$

$$X = u$$

$$Y = 0$$

$$T = 0$$

DOUBLE PENDULUM CART (continued)

M_5 : horizontal free motion

variables:

$$x, y, \theta, X, Y, T$$

behavioral equations:

$$\begin{aligned}y &= 0 \\ \theta &= -\frac{\pi}{2} \\ X &= 0 \\ T &= 0\end{aligned}$$

DOUBLE PENDULUM CART (continued)

Write the interconnection laws:

For each interconnection, the 2-D mechanical interconnection laws:

$$x' = x''$$

$$y' = y''$$

$$\theta' = -\theta''$$

$$X' = -X''$$

$$Y' = -Y''$$

$$T' = T''$$

etc., etc.

DOUBLE PENDULUM CART (continued)

After suitable manipulations, we obtain:

ELIMINATION

Consider

$$R\left(\frac{d}{dt}\right)w = M\left(\frac{d}{dt}\right)\ell \quad (\star)$$

w : w -dimensional, ℓ : 1-dimensional

$$\begin{aligned} R &\in \mathbb{R}^{\bullet \times w}[\xi], \quad M \in \mathbb{R}^{\bullet \times 1}[\xi] \\ R(\xi) &= R_0 + R_1\xi \cdots R_n\xi^n \\ M(\xi) &= M_0 + M_1\xi \cdots M_n\xi^n \end{aligned}$$

.....

Define

$$\mathfrak{B}_{\text{full}} = \{w : \mathbb{R} \rightarrow \mathbb{R}^w \times \mathbb{R}^1 \mid (\star) \text{ holds}\}$$

$$\mathfrak{B} = \{w \mid \text{there is an } \ell \text{ such that } (\star) \text{ holds}\}$$

ELIMINATION (continued)

Questions:

- Is \mathfrak{B} described by a constant-coefficient differential equation?

i.e., is there $R' \in \mathbb{R}^{\bullet \times w}[\xi]$ such that

$$R' \left(\frac{d}{dt} \right) = 0 \quad R' \in \mathbb{R}^{\bullet \times w}[\xi]$$

is a kernel representation of \mathfrak{B} ?

- If so, find an algorithm

$$(R, M) \mapsto R'$$

- Nonlinear, PDE generalizations?

ELIMINATION (continued)

Theorem:

There indeed exists an R' such that

\mathfrak{B} is the solution set of

$$R' \left(\frac{d}{dt} \right) w = 0$$

.....

Projection of linear differential
behavior = same!

.....

Algorithm for computing R' :

$n \in \mathbb{R}^{1 \times \bullet}$ is *annihilator* for M if
 $nM = 0$.

Set of annihilators = module;
finitely generated

Generators = (n_1, n_2, \dots, n_g) ;

defines matrix N

$R' = NR$.

To be made into computer algebra

Pseudocode for elimination

```
function  $R'$ =LATELIM( $R, M$ );  
% Input : Polynomial matrices  $M$  and  
 $R$  with same number of rows  
% Output: Polynomial matrix  $R'$  such  
that  $R'(\frac{d}{dt})w = 0$  is the manifest  
% behavior corresponding to the  
hybrid representation  
%  $R(\frac{d}{dt})w = M(\frac{d}{dt})\ell$   
%  
% At each step, the degree of one of  
thw rows of  $M$  is decreased by adding  
a  
% suitable polynomial combination of  
other rows. The same operation  
% is applied to the rows of  $R$ . When  
a row of  $M$  becomes  $0$ , the  
% corresponding row of  $R$  becomes part  
of  $R'$ .
```

```

[M, R]=Order(M, R);
% order rows so that rows of M
are in decreasing degree order
i=rowdim(M); p=i;
k=1;

while(all(M(i,:) == 0)) %check
if M has zero rows
    R'[k,:]=R[i,:]; %corresponding
rows of R are in K
    i=i-1; p=p-1; k=
k+1;
end

while ((i ≥ 1)) %reduction procedure

    if (∃ real n s.t. Mhc(i,:) =
nMhc(i+1:p,:)) then
        % check if highest coefficient
is linearly dependent
        % from those of rows of lower
or equal degree

        h=polann(n, M); %polynomial
vector to reduce degree
        m=M(i,:)-hM(i+1:
p,:); % degree(m)< degree(M(i,:))
        r=R(i,:)-hR(i+1:p,:
); %same combination on R

```



```

         $[M, R] = \text{Eliminate}(M, R, i);$ 
% eliminate row

        if ( $m \neq 0$ ) then
             $[M, R, j] = \text{Insert}(M, R, m, r);$ 
            % insert  $m$  and  $r$  as  $j$ -th
row of  $M$  and  $R$ 
            % so as to keep rows of
 $M$  ordered by degree

             $i = j;$ 

            else % the combination of
rows of  $M$  is zero
                 $R'[k, :] = r$  %combination
of rows of  $R$  in  $K$ 
                 $i = i - 1; \quad p = p -$ 
1;  $k = k + 1;$ 
                end

            else  $i = i - 1;$ 
                end

        end
end

```

CONTROLLABILITY

The time-invariant system $(\mathbb{R}, \mathbb{W}, \mathfrak{B})$
is said to be

controllable

if for all $w_1 \in \mathfrak{B}, w_2 \in \mathfrak{B}$,

there exists $t' \geq 0$ and
 $w \in \mathfrak{B}$

such that

$$w(t) = \begin{cases} w_1(t) & \text{for } t < 0 \\ w_2(t - t') & \text{for } t \geq t' \end{cases}$$

CONTROLLABILITY (continued)

Questions:

- Is the system described by

$$R\left(\frac{d}{dt}\right)w = 0$$

controllable?

- Find an effective algorithm for verifying controllability in terms of R_0, R_1, \dots, R_n where

$$R(\xi) = R_0 + R_1\xi + \dots + R_n\xi^n$$

- Nonlinear, PDE generalizations?
- Application/relevance in control

CONTROLLABILITY (continued)

Theorem:

The following are equivalent

1. $R(\frac{d}{dt})w = 0$ defines a controllable system
2. $\text{rank}(R(\lambda))$ is independent of λ for $\lambda \in \mathbb{C}$
3. The behavior \mathfrak{B} is the manifest behavior of

$$w = M\left(\frac{d}{dt}\right)\ell$$

for some $M \in \mathbb{R}^{w \times \bullet}$

\mathfrak{B} admits an “*image representation*”.

CONTROLLABILITY Verification

Idea

Given R

$f \in \mathbb{R}^{\bullet \times 1}[\xi]$ belongs to the
SYZGY of R if $Rf = 0$

SYZGY = module, finitely gen.
generators $\{r'_1, r'_2, \dots, r'_g\}$
form matrix R'

annihilators of $R =: \mathcal{N}_R$
annihilators of $R' =: \mathcal{N}_{R'}$

Controllability test:

$$\mathcal{N}_R = \mathcal{N}_{R'}?$$

→ Computer algebra

Pseudocode for verifying controllability

```
function [ $M$ ,obs]=RPR( $D$ );  
% Input :  $p \times \ell$  polynomial matrix  
 $D$   
% Output: boolean variable obs,  $p \times$   
 $\ell$  polynomial matrix  $M$   
%  
% Builds 0-degree rows generated by  
the rows of  $M$   
% and checks if they have rank= $\ell$ ; if  
so, output variable obs=1.  
% At each step the degree of one row  
is decreased by adding a suitable  
% polynomial combination of other rows.  
This goes on until no lowering  
% is possible, or the above condition  
is satisfied.  
% Output matrix  $M$  is the result of  
such reduction.  
%  
%  
 $M = D$   
 $M$ =Order( $M$ ); % order rows in decreasing
```

```

degree order
    obs=(rank( $M^0$ ) ==  $\ell$ ); % check
if 0-degree rows already enough
     $p$ =rowdim( $M$ );
     $i$ = $p$ -rowdim( $M^0$ ); % otherwise start
from first row of higher degree

    while ((not obs) and ( $i \geq 1$ ))
        % Reduce until enough 0-degree rows
or no more reduction possible

        if ( $\exists$  real  $n$  s.t.  $M_{hc}(i,:) =$ 
 $nM_{hc}(i+1:p,:)$ ) then
            % check if highest coefficient
is linearly dependent
            % from those of rows of lower
or equal degree

             $h$ =polann( $n, M$ ); %polynomial
vector to reduce degree
             $m = M(i,:) - hM(i+1 :$ 
 $p,:)$ ; % degree( $m$ )< degree( $M(i,:)$ )
             $M$ =Eliminate( $M, i$ ); % eliminate
row

            if ( $m \neq 0$ ) then
                [ $M, j$ ]=Insert( $M, m$ );
                % if new vector is not
0 insert it as  $j$ -th row
                % so as to keep rows ordered

```

```

by degree
    if (degree( $m$ ) == 0)
then
    obs=(rank( $M^0$ ) ==
 $\ell$ );
         $i = j - 1$ ; % determine
new row to examine
        else  $i = j$  end
        else  $p = p - 1$ ;  $i = i - 1$ ;
end
    else  $i = i - 1$ ;
    end
end

```

```

function ctr=CTRB(R);

%
% Input: Polynomial matrix R with


p rows and q columns


% Output: Boolean variable ctr=1 if
ker( $R(\frac{d}{dt})$ ) is controllable
%
%
    [P,ctr]=RPR(RT);
    % Check if the 0-degree columns
generated by
    % the columns of R have rank p
if (not ctr) then
     $P = P^T$ ;
     $P=COLPRP(P)$ ; % Bring P in
column proper form
    if (rowdim(P)>coldim(P)) then
        % if the matrix was not of full
row rank

        [P,ctr]=RPR(P);
        %Check if the 0=degree rows
generated by
        % column proper form have
rank=q
    end
end
end

```


REPRESENTATIONS

$$R\left(\frac{d}{dt}\right)w = 0$$

$$w = M\left(\frac{d}{dt}\right)\ell$$

$$R\left(\frac{d}{dt}\right)w = M\left(\frac{d}{dt}\right)\ell$$

$$E\frac{d}{dt}x + Fx + Gw = 0$$

$$\begin{aligned}\frac{d}{dt}x &= Ax + Bu \\ y &= Cx + Du \\ w &\approx (u, y)\end{aligned}$$

$$\begin{aligned}P\left(\frac{d}{dt}\right)y &= Q\left(\frac{d}{dt}\right)u \\ w &\approx (u, y)\end{aligned}$$

$$\begin{aligned}y &= G(s)u \\ w &\approx (u, y)\end{aligned}$$

! Algorithms for passing among them
! Algorithms for testing various prop-
erties ! Algorithms for synthesis

SIMULATION

$$R\left(\frac{d}{dt}\right)w = 0$$

will, of course, have many solutions,
due to

- free variables among (w_1, w_2, \dots, w_m)
- free initial conditions

In order to simulate a response, we
need additional data

$$K\left(\frac{d}{dt}\right)w = f \quad f : \mathbb{R} \rightarrow \mathbb{R}^\bullet \text{ given}$$

$$S\left(\frac{d}{dt}\right)w(0) = a \quad a \in \mathbb{R}^\bullet \text{ given}$$

Does there exist a solution?

Does there exist a unique solution?

If so, algorithm

$$(R, K, S, f, a) \mapsto w$$

full plant behavior)

$$\mathcal{P}_{\text{full}} = \{(v, c) \in \mathcal{C}^\infty(\mathbb{R}, \mathbb{R}^{v+c}) \mid (v, c) \text{ satisfies the plant}$$

plant behavior

$$\mathcal{P} = \{v \in \mathcal{C}^\infty(\mathbb{R}, \mathbb{R}^v) \mid \exists c \in \mathcal{C}^\infty(\mathbb{R}, \mathbb{R}^c) \text{ such that } (v, c)$$

controlled behavior \mathcal{K} defined by

$$\mathcal{K} = \{v \in \mathcal{C}^\infty(\mathbb{R}, \mathbb{R}^v) \mid \exists c \in \mathcal{C} \text{ such that } (v, c) \in \mathcal{P}_{\text{full}}\}.$$

For what $\mathcal{K} \in \mathcal{L}^v$ does there exist a $\mathcal{C} \in \mathcal{L}^c$ that implements \mathcal{K} ?

\mathcal{C} implements \mathcal{K} if the above relation holds between \mathcal{C} and \mathcal{K} .

For what $\mathcal{K} \in \mathcal{L}^v$ does there exist a $\mathcal{C} \in \mathcal{L}^c$ that implements \mathcal{K} ?

Sheet

hidden behavior and is denoted as

\mathcal{N} . It is formally defined as

$$\mathcal{N} = \{v \in \mathcal{P} \mid (v, 0) \in \mathcal{P}_{\text{full}}\}.$$

Theorem 1 (Controller implementability theorem)
: *Let $\mathcal{P}_{\text{full}} \in \mathcal{L}^{v+c}$ be the full plant behavior, $\mathcal{P} \in \mathcal{L}^v$ the manifest plant behavior, and \mathcal{N} the hidden behavior. Then $\mathcal{K} \in \mathcal{L}^v$ is implementable by a controller $\mathcal{C} \in \mathcal{L}^c$ acting on the control variables if and only if*

$$\mathcal{N} \subset \mathcal{K} \subset \mathcal{P}.$$

Sheet Assume also that the transfer function $G_{(d,u)\mapsto(f,y)}$ associated with $\mathcal{P}_{\text{full}}$ has the following properties:

- (i) $G_{(d,u)\mapsto(f,y)}$ is proper,
- (ii) $G_{u\mapsto f}^{\infty}$ is injective,
- (iii) $G_{d\mapsto y}^{\infty}$ is surjective, and
- (iv) $G_{u\mapsto y}^{\infty} = 0$.

Let $\mathcal{N} \in \mathfrak{L}^{d+f}$ be the hidden behavior, and $\mathcal{P} \in \mathfrak{L}^{d+f}$ be the plant behavior associated with $\mathcal{P}_{\text{full}}$. Assume that the behavior $\mathcal{K} \in \mathfrak{L}^{d+f}$ satisfies:

- (v) $\mathcal{N} \subset \mathcal{K} \subset \mathcal{P}$, i.e., \mathcal{K} is an implementable controlled behavior,
- (vi) in \mathcal{K} , d is input and f is output, and
- (vii) the transfer function $K_{d\mapsto f}$ from d to f in \mathcal{K} is proper.

Then there exists a controller $\mathcal{C} \in \mathfrak{L}^{u+y}$ such that

1. \mathcal{C} implements \mathcal{K} ,
2. in \mathcal{C} , y is input and u is output, and
3. the transfer function $C_{y\mapsto u}$ from y to u in \mathcal{C} is proper.

SALIENT FEATURES

- Dynamical system = a behavior input/output structure: important special case
- First principles models \implies latent variables
state variables: important special case
- Control = interconnection feedback : important special case
- Modelling complex systems = tearing & zooming
input-to-output cascade and feedback: limited special case

terminal	equations
electrical	$V_1 = V_2, I_1 + I_2 = 0$
1-D mechanical	$q_1 - q_2, F_1 + F_2 = 0$
2-D mechanical	$x_1 = x_2, y_1 = y_2, \theta_1 = -\theta_2,$ $X_1 + X_2 = 0, Y_1 + Y_2 = 0, T_1 = T_2,$
thermal	$T_1 = T_2, Q_1 + Q_2 = 0$
fluidic	$p_1 = p_2, f_1 + f_2 = 0$
logical output \rightarrow input	$u=y$
etc. etc.	etc. etc.

RLC-circuit

Problem: Model the relation between
 (V_1, I_1, V_2, I_2)

STABILITY

The system

$$R\left(\frac{d}{dt}\right)w = 0$$

is *stable* if $(w \in \mathfrak{B}) \Rightarrow (w(t) \rightarrow 0 \text{ for } t \rightarrow \infty)$

Stable \Leftrightarrow

$(\lambda \in \mathbb{C} \text{ and } \text{rank}(R(\lambda)) < \bar{w}) \Rightarrow (R_\varepsilon(\lambda) < 0)$

is *stabilizable* if for all $w \in \mathfrak{B}$ there is $w' \in \mathfrak{B}$ such that

$$\begin{aligned} w'(t) &= w(t) && \text{for } t < 0 \\ w'(t) &\rightarrow 0 && \text{for } t \rightarrow \infty \end{aligned}$$

Stabilizable \Leftrightarrow

$(\lambda \in \mathbb{C} \text{ and } \text{rank}(R(\lambda)) < \text{rank}(R)) \Rightarrow (R_\varepsilon(\lambda) < 0)$

STABILIZABILITY

A plant is stabilizable by a regular control interconnection if and only if

1. \mathcal{N} is stable
2. \mathcal{P} is stabilizable

Note : \mathcal{N} is stable = “detectability”
i.e $c = 0 \Rightarrow w(t) \rightarrow 0$ for $t \rightarrow \infty$